# A Clustering Expert with Bandit Feedback of their Performance in Many Tasks

Maximilian Graf and Victor Thuot

Group serminar: Mathematical Statistics.
Potsdam, 22<sup>rd</sup> November 2024
University of Potsdam

#### **Model: Performance Matrix**

Consider 
$$\Omega$$
 "experts"

Consider  $\Omega$  "tadas"

The performance of experts  $i \in \{1,...,n\}$  on the task  $j \in \{1,...,l\}$  is given by this  $i \in \mathbb{R}$ .

Performance:  $M = \{M, 1, ..., M, d\}$ 

That is  $M = \{M, 1, ..., M, d\}$ 

That is  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ..., M, d\}$ 

The performance  $M = \{M, 1, ...$ 

01

#### **Model: Performance Matrix**

Consider n "experts"
Consider d'tarby The performance of experts  $i \in \{1, ..., n\}$  on the task  $j \in \{1, ..., l\}$  is given by  $Mij \in \mathbb{R}$ .  $Mi \in \mathcal{L}$   $Mi,1 \cdot Mi,d \leftarrow \text{expect } i$   $Mi,1 \cdot Mn,d$ 

# Model: Sequential and adaptive learning protocol

reguential: one data at a time adaptive: the choice (It, It) is based on the passed

## Model: Sequential and adaptive learning protocol

sequential: one data at a time

adaptive: the choice (It, It) is based on the passed

# Model: Stationary and subGaussian feedback

Assumption 1 (stationasity)
There exists 
$$n \times d$$
 distributions on  $R: \mathcal{D}_{i,j}$ 

L( $X_t \mid T_{n,T_n}, X_{n,---}, T_{t,n,T_{t-1}}, X_{t-1}, T_{t,n,T_{t-1}} = X_{t,n,T_{t-1}}$ 

feedback.

Passed dicisions decisions at time t

Exp  $X_{i,j}$ 

Assumption 2 (subGaussian noise)

For all  $(i,j)$ , if  $X_i \sim \mathcal{D}_{i,j}$ 

then  $X_{-}\mathcal{U}_{i,j}$  is 1-sub-Gaussian

# Model: Hidden partition

Assumption 1: assume that these exists {  $\mu^b \in \mathbb{R}^d$ such that, for any expert ie f1,..., ng, Mi E of Ma, Mb &  $M = \begin{pmatrix} M_{1,1} & \dots & M_{n,d} \\ \vdots & \vdots & \vdots \\ M_{i,1} & \dots & M_{i,d} \\ \vdots & \vdots & \vdots \\ M_{n,1} & \dots & M_{n,d} \end{pmatrix}$   $M_{n,1} & \dots & M_{n,d}$ 

# Model: Hidden partition

Assumption 1: assume that these exists {  $\mu^b \in \mathbb{R}^d$ 

such that, for any expert is f1,..., ng, Mi E of Ma, Mb

Notation

# Model: Hidden partition

Assumption 1: assume that these exists {  $\mu^b \in \mathbb{R}^d$ 

such that, for any expert ie f1,..., ng, Mi E of Ma, Mb &

Noration

## Model: Sequential and adaptive learning protocol

> sampling rule: (Tt, Tt) Es 1

## Model: Sequential and adaptive learning protocol

# Objective: PAC setting

First objective

Given  $S \in (0,1)$ , an algorithm of is S-PAC if.  $PA, S = S^a \text{ or } \hat{S} = S^b > 1-S$ 

# Objective: PAC setting

First objective Given SE(0,1), an algorithm et is S-PAC if.  $P_{A, \gamma}(\hat{S} = S^{\circ} \circ \hat{S} = S^{\circ}) > 1 - S$ Should recover 5°1156 proba induced by elgosithm of up to permutation and environment )

# Objective: PAC setting

First objective

Given SE(0,1), an algorithm et is S-PAC if.

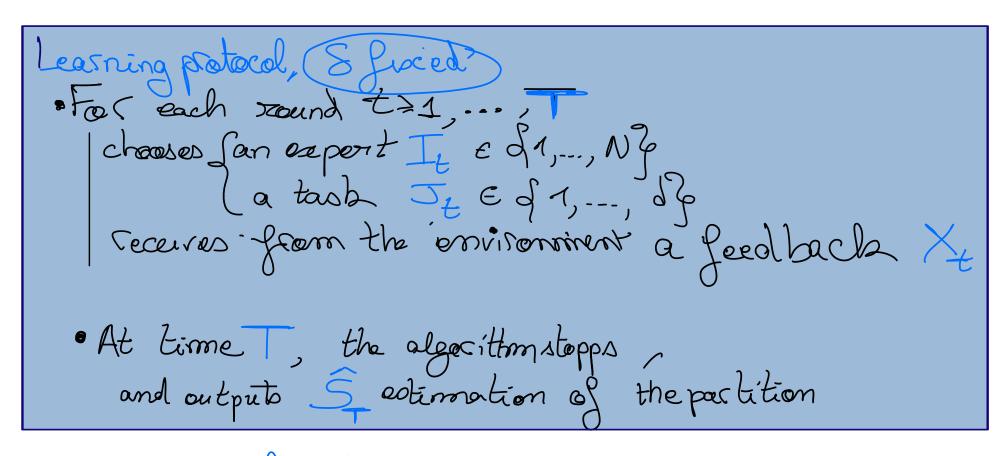
 $P_{A,y}(\hat{S}_{z}, \hat{S}_{z}, \hat{S}_{$ 

Second objective

Construct of with T as small as possible while maintaining \*

Ry Tis a stepping time designed

# Model: Sequential and adaptive learning protocol



- Decision rule: (It, It) Es 1
- Destopping rule: Telepping time.

Remaques:

main objective:

balance S exploration to Cearn the gap vector & (exploration to classify the arms (exploration 12)

D'Cinks with (adaptive) signal detection\_

I link to the previous paper

#### Lower bound:

# Theorem For any algorithm of, S-PAC for the clustering problem, there exists $\sigma, \tau$ two permutations such that, if $M_{s,\tau}$ is constructed withe M by permutation $(\sigma, \tau)$ , then Pt. $M_{s,\tau}$ $\left( T \ge \frac{2d}{\theta \, \text{II} \, \text{Syz}} \, \log \left( \frac{1}{68} \right) \right) \ge S$

where 
$$\Delta = \mu^a - \mu^b \in \mathbb{R}^d$$
 is the gap exector

### Lower bound: intuition

imagine that  $M^{9}=(9,--,0)$  $\mu = (\mu, --, \mu, 0, ---, 0)$  and  $|S^b| = \Theta n$ · 112/2 = 2/12 in order to detect one expert in S, prohability on expert to in St. \* | que need to sample far each expert de toubs to we need 1/2 top (1) souples from an (expert, took)-comple Le decide reether the entry of Mis Oct 11

# Algorithm: three-steps procedure

Main structure of the algorithm: <u>Step 1</u> : idenfify one expert in each group Sta∈Sa Tib∈Sb Step 2: cellect information on the structure of  $\Delta$  (gap vector  $M^a - \mu^b$ ) and choose a task  $\mathcal{J} \in \mathcal{J}_{1,-}, \mathcal{J}_{p}$  such that  $|\Delta_{\mathcal{J}}|$  is large Step 3: clarify each expest based on its perferenance on the task j

Step 1

L

Signal detection

Step 2 2 - Best erm identification (and variants) Step 3

Livery classification

# Sequential Halving: algorithm

« Sequentiel Holking is a procedure for BAIT which has guaranties for: ≥ ε- bestasm identification ⇒ (ε,m) - BAI

If we consider a b-wered bandit, with asms of an, ---, and

budget T, S. Jan, --, able set of arms

For u = 0, ..., Plage (b):

sample I times each arm in Su

Plage(b) #Su

· Su+1 = half boot arm from Su

# Sequential Halving: guaranties

bourns reith means [h,---, lk] ordered as his,---, lb)

E-BAI

(E,m)-BAI

Remark: we have to adapt the procedure as:

—> we have positive and regative entries.

—> our objective is signal detection and not BAT

# Step 1: representatives identification ( \*\*expest)

Objective: identify two experts to fa, Fb & \$1,..., n} such that read is are in different groups w.h.p. (>1-8)

Algo P Sequential Halving

Deubling trich become Dand D are unknown

Deubling: Stronly on a subset of expert-task couples

Trough two sample testing

Garanties  $\Rightarrow$  algo which identifies  $\Re_a$ ,  $\Im_b$  which are generated to be representatives supposed budget (in A-8) - quantile)  $\frac{1}{2} \frac{\partial^2 A}{\partial a^2} \frac$ 

# Step 2: learning the gap vector

Balance the budget used to classify the rasms. With step 1, we have access to Staes expert with performance ree Rd Libes by who IRd In particular, we have to identify a task 3G1,-, of such that Is is large.

Again rue use SH and we are interested in (E, m) - BAI

# Step 3: classification

· We know Sta ∈ S° respect with perforemence us

The ∈ S° respect with perforemence us

The est took such that | Aj| large.

There is each expect on took of

There is such as performence us

There is took such that | Aj| large.

## Upper bound

Theorem The three-steps procedure is S-PAC and with probability at least 1-8, it holds that  $\frac{d}{d} \log \left(\frac{1}{8}\right) + \min \left(\frac{1}{8}\right) + \frac{1}{160} \log \left(\frac{1}{8}\right) + \frac{1}{160} \log$ 

where  $N_{(n)} \geq N_{(n)} > --- \geq N_{(n)}$  are the entries of  $\Delta$  ordered by absolute value.

#### Lower bound

Theorem

For any S-PAC algorithm, there exists  $s, \tau$  permutations of the rows and columns, such that  $\frac{d}{dt} = \frac{d}{dt} = \frac{d}$ 

where  $M_{(n)}|_{\geq}M_{(n)}|_{\geq}---\geq |M_{(n)}|$  are the entries of  $\Delta$  ordered by absolute value.

# Finding Candidates

- · assume & was known and  $\Delta = \mu^a \mu^b$ was S Sparse such that every mon-zero entry is h > 0
- Sampling  $Ci_1, j_1), ..., Ci_{b}, j_{b}) \sim \mathcal{M}([N] \times [d])$ with  $\phi \gtrsim \frac{d}{s \oplus log}(\frac{1}{s})$  yields

 $\Delta_{ikjk} = h$  for some  $k = 1, ..., \phi$ 

- · halog (f) samples required to decide

  Mij=0 or Mij=h correct W. P. 31-5
- · finding (inin) requires de la log)

# Finding Candidates

- in general,  $\Delta$  is not S-sparse, but  $\exists s \in [d]$  such that  $\|\Delta\|_{2}^{2} \leq \log(2d) s \cdot \Delta_{CS}$ ,  $\|\Delta\|_{2}^{2} \leq \log(2d) s \cdot \Delta_{CS}$ ,  $\|\Delta\|_{2}^{2} \leq \log(2d) s \cdot \Delta_{CS}$ ,  $\|\Delta\|_{2}^{2} \leq \log(2d) s \cdot \Delta_{CS}$
- o still:  $\|A\|_{2}^{2}$ , s and  $\Theta$  unknown

  adaptivity by doubling trick of

# Finding camdidates

• again: assume  $\|\Delta\|_2^{\ell}$ , s and  $\Theta$  were known consider  $h = \frac{\|\Delta\|_2^2}{s \log(2d)}$  iid
• recall: Sampling Cin,  $j_1$ , ..., Cip,  $j_2$ )  $\sim \mathcal{M}([N] \times [d])$  with  $\phi \gtrsim \frac{d}{s\Theta} \log(\frac{1}{s})$  yields  $\Delta i_k j_k \geq h$  for  $c \cdot log(f) k = 1, ..., \phi$ · finding one entry with 1 dixix 1 > \frac{h}{2} would require a total budget  $\sim can we$   $\gtrsim \phi \cdot \frac{\log(\phi)}{h^2} \gtrsim \frac{d}{\|A\|_{\ell}^2} \log(\frac{1}{\sigma})^2 do bette?$ up to polylog in d and N

Halving (intuition) Sequential \$ ~ = log ( = ) (in ight) from which o recall: Dinin≥h for at least c. (og (-1) arms · idea: - Sample I times each arm (ii) - obtain dij as averages and das corresponding median - remove all arms with sij < \$\D - T = 2. T, repeat until only one arm left o algorithm terminates after L= [loga(\$)] steps · Se set after I halving steps, Ue={ci,j)es: 1/10/> (1-21)h3

· We show:  $\frac{|u_{i}|}{|S_{i}|}$  is mon-decreasing  $u.h.p. \Rightarrow S_{1} + \emptyset$ 

Case 1: 1 < 1 (he+ + he)  $h_{i} := \left(1 - \frac{k}{2L}\right)h$ hen I at least half of the remains in Se+1 Case 2: 1 > \frac{1}{2}(hen + he) a) number of ciij) ese with Milkheth large -> at least half of them not in Sex b) it small -> Ules automatically large enough