

Active clustering with bandit feedback

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- 1 Setting: Active Clustering Problem
- 2 Contributions: Gaussian setting
- 3 Information-theoretic Lower bound
- 4 Upper Bound: the ACB Algorithm

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- the algorithm receives X_t , s.t.,
conditionally on $A_t = a$, $X_t \sim \nu_a$

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- $(\nu_1, \dots, \nu_N)$: **environment**

Hidden partition assumption

The **Active Clustering Problem (ACP)** ([Yang et al., 2024]) is defined by

- a bandit environment with N arms that have to be partitioned
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Assumption 1: hidden partition

There exists a **hidden partition** G^* of $[N]$ into K groups.
For $a, b \in [N]$, a, b are in the **same group** $\Leftrightarrow \mu_a = \mu_b$.

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- each group $k \in [K] \leftrightarrow$ center $\mu(k)$
- The k -th group, is $G_k^* := \{a \in [N]; \mu_a = \mu(k)\}$
- assume the groups nonempty and disjoint
- K is known (and also N and d)

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δ -PAC algorithm

Given $\delta \in (0, 1)$, an algorithm π for the ACP is said to be δ -PAC on a collection of environments $\mathcal{E}$ if for any $\nu = (\nu_1, \dots, \nu_N) \in \mathcal{E}$, then

$$\mathbb{P}_{\pi, \nu}(\hat{G} \sim G^*) \geq 1 - \delta .$$

where $\sim$ means that the partition is exact up to permutation of the groups.

Optimal worst case budget

We define the **optimal worst case (expected) budget** on the collection $\mathcal{E}$ as

$$T^*(\delta, \mathcal{E}) = \inf_{\pi} \sup_{\nu \in \mathcal{E}} \mathbb{E}_{\pi, \nu}[\tau] ,$$

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Gaussian noise

Assumption 2: Gaussian noise

- 1 there exists Σ_a s.t. $\nu_a = \mathcal{N}(\mu_a, \Sigma_a)$
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- 2 we know σ s.t., $\sigma^2 \geq \sup_{a \in [N]} \|\Sigma_a\|_{op}$

Parameters of interest

First, we consider $\Delta_* = \Delta_*(\nu)$ for the **minimal gap**,

$$\Delta_* = \min_{k \neq k'} \|\mu(k) - \mu(k')\| > 0 .$$

Besides, we denote m_* as the **size of the smallest group**,

$$m_* = \min_{k \in [K]} |G_k^*| .$$

Construction of a family of environments

- Consider $\Delta > 0$, and $m \geq 1$, we define

$$\mathcal{E} := \mathcal{E}(\Delta, m) ,$$

as the collection of environments ν such that $\Delta_*(\nu) \geq \Delta$,
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- We study the optimal worst case budget on $\mathcal{E}(\Delta, m)$

$$T^*(\delta, \mathcal{E}(\Delta, m)) .$$

Benchmark: the uniform sampling strategy

Uniform sampling strategy (US):

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$$\tau = NT \gtrsim \frac{\sigma^2}{\Delta^2} \left[N(\log(N) \vee K) + \sqrt{dKN(\log(N) \vee K)} \right]$$

([Royer, 2017, Giraud and Verzelen, 2019,
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- Information-computation gap

Benchmark: optimal result in the asymptotic regime $\delta \rightarrow 0$

The problem was introduced in [Yang et al., 2024], they provide:

- 1 Instance-dependent lower bound
- 2 An algorithm optimal in the asymptotic regime $\delta \rightarrow 0$
 - for instance, with $m = N/K$, equal distance between the groups, and equal variance $\sigma^2 I_d$,

$$\lim_{\delta \rightarrow 0} \frac{T^*(\delta, \mathcal{E}(\Delta, m))}{\log(1/\delta)} = 2 \frac{\sigma^2}{\Delta^2} (N + K)$$

- methodology: ([Garivier and Kaufmann, 2016])

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Yes,
$$T^* \simeq \frac{\sigma^2}{\Delta^2} \left[N \log \left(\frac{N}{\delta} \right) + \sqrt{dKN \log \left(\frac{N}{\delta} \right)} \right].$$

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- 3 Is there an information-computation gap for ACP?

No, there is no computational gap, ACB is optimal in most emblematic regimes (e.g., balanced setting).

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Information-theoretic lower bound

Theorem 1

There exists a universal constant $c > 0$, such that the following holds for any $\Delta > 0$, any $m \geq 2$, any $\delta \in (0, 1/12)$, and any $N \geq mK$:

$$T^*(\delta, \mathcal{E}(\Delta, m)) \geq c \left[\frac{\sigma^2}{\Delta^2} N \log \left(\frac{N}{\delta} \right) + \frac{\sigma^2}{\Delta^2} \sqrt{dKN \log \left(\frac{N}{\delta} \right)} \right] .$$

Comments

Dichotomy between low-dimensional and high-dimensional clustering problems:

$$T^* \geq c \left[\frac{\sigma^2}{\Delta^2} N \log \left(\frac{N}{\delta} \right) + \frac{\sigma^2}{\Delta^2} \sqrt{dKN \log \left(\frac{N}{\delta} \right)} \right] .$$

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- In low-dimension, the bound is obtained by reducing the problem to a specific instance of the thresholding bandit problem ([Cheshire et al., 2020, Chen and Li, 2015, Chen et al., 2014]).
- Simpler problem: $K = 2, d = 1, \mu_a \in \{0, \Delta\}$ with Δ known

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- In high-dimension, we use a series of reduction, driven by two ideas:
 - 1 Bayesian approach: chose a "good" prior on the unknown centers
 - 2 ACP is at least as "difficult", as the (active) supervised problem where we would know the group of all the arms except one

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Upper bound: the ACB algorithm

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Main structure of the algorithm:

- 1 identify $\hat{S}$ a set of K arms with exactly one arm from each group
- 2 estimate the unknown centres of the groups using $\hat{S}$
- 3 sample uniformly the remaining arms and use a distance-based classifier

First step: identification of representatives

Imagine that $\hat{S}$ contains $k < K$ arms from different clusters. Until $|\hat{S}| < K$,

- take randomly a new candidate $b \in [N]$
- perform a sequence of tests to decide whether b should be added to $\hat{S}$ or not
- the objective is to reject quickly arms whose groups are already represented

We use for the tests [sub-sampling](#) and high dimensional [two-sample testing](#).

Similar to an elimination technique

[de Heide et al., 2021, Jamieson and Nowak, 2014].

Upper bound: main theorem

Theorem 2

Let $\delta > 0$. The ACB algorithm is δ -PAC on $\mathcal{E}(\Delta, m)$.
If we assume that $m \geq \log(K)$, then

$$\mathbb{E}_{ACB, \nu}[\tau] \leq c \frac{\sigma^2}{\Delta^2} \left[N \log(N/\delta) + \sqrt{dNK \log(N/\delta)} + \sqrt{d} \frac{N \log(K)}{m} \right]$$

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- c is independent of all parameters
- we also give a bound with high probability

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$$\left(m \geq \sqrt{\frac{N}{K}} \frac{\log(K)}{\sqrt{\log(N/\delta)}} \right)$$

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- 2 the upper bound is optimal if m is large enough,
 $\left(m \geq \sqrt{\frac{N}{K}} \frac{\log(K)}{\sqrt{\log(N/\delta)}} \right)$
- 3 the algorithm is polynomial $\rightarrow$ no information-computation gap

Take home message

- 1 There is no computational gap for the ACP.
- 2 We provide a lower bound on the budget

$$T^* \geq c \left[\frac{\sigma^2}{\Delta^2} N \log \left(\frac{N}{\delta} \right) + \frac{\sigma^2}{\Delta^2} \sqrt{dKN \log \left(\frac{N}{\delta} \right)} \right] .$$

- 3 We provide a polynomial time δ -PAC algorithm called ACB, together with an upper bound on its budget which matches the lower bound for mild assumption on m_* .

Open questions

- 1 Design an environment-dependent non-asymptotic bound
- 2 Derive other active vs batch comparison

Thank you !

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- Exemples : **bounded noise or Gaussian noise.**

First step: identification of representatives

Consider for $s \in [r]$ the s -th test performed on the candidate b .
For $a \in \hat{S}$, we compute the statistic

$$\langle \bar{\mu}_{b,s} - \hat{\mu}_a, \bar{\mu}'_{b,s} - \hat{\mu}'_a \rangle$$

- $\hat{\mu}_{b,s}, \hat{\mu}'_{b,s}$ are two independent estimation of μ_b computed with n_s samples
- $\hat{\mu}_a, \hat{\mu}'_a$ are estimates of μ_a computed with $n_{\max}$ samples
- the expectation of this statistic is $\|\mu_a - \mu_b\|^2$
- we reject b if it is smaller than $\Delta^2/2$ for some $a \in \hat{S}$.

We use sub-Gaussian concentration to choose the tuning parameters $n_s, n_{\max}, r$.

Second step: classification

Imagine that $\hat{S} = \{a_1, \dots, a_K\}$ contains one arm from each group.

- 1 First, for $j \in [K]$, label a_j with j and estimate $\mu(j)$ with two independent means using $2J$ samples.
- 2 Then, for each $b \in [N] \setminus \hat{S}$, labels b in the group

$$\operatorname{argmin}_{j=1,\dots,K} \left\langle \hat{\mu}_b - \hat{\mu}(j), \hat{\mu}'_b - \hat{\mu}'(j) \right\rangle$$

$\hat{\mu}_b, \hat{\mu}'_b$ are computed with $I = KJ/N$ samples